# Database tools for the large scale computation of Maximally Mutable Laurent polynomials 

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## Motivation

## Context

Fanosearch Project: study Fano varieties using Mirror Symmetry. [Coates, Corti, Galkin, Golyshev, Kasprzyk (arXiv:1212.1722)]

- Two sides of the mirror:
$\left\{\begin{array}{c}\text { some Laurent } \\ \text { polynomials }\end{array}\right\} /\{$ mutations $\} \Leftrightarrow\left\{\begin{array}{c}\text { Fano } \\ \text { varieties }\end{array}\right\} /\{$ deformations $\}$
- Laurent polynomials are simpler than Fano manifolds.


## Motivation

## Definition

Fix a lattice $N \equiv \mathbb{Z}^{d}$.
A canonical Fano polytope $P \subset N$ is a convex lattice polytope such that:

- the vertices $\mathcal{V}(P)$ are integral lattice points.
- The only lattice point in $\stackrel{\circ}{ }$ is the origin of the lattice.
- $\operatorname{dim}(P)=d$.


## A special case

- If the polar polytope $P^{*}$ is itself an integral lattice polytope, we call $P$ reflexive.


## Motivation

"Interesting" Laurent polynomials $f$ :

- Given a canonical Fano polytope $P$,

$$
\text { Newt } f=P, \quad f=\sum_{\gamma \in N \cap P} a_{\gamma} x^{\gamma}
$$

## Definition

Given a vector $w \in \operatorname{Hom}(N, \mathbb{Z})$, and $u \in \mathbb{C}\left[w^{\perp}\right]$, a mutation $\Phi: \mathbb{C}[N] \rightarrow \mathbb{C}[N]$ defined by:

$$
\begin{gathered}
\Phi: x^{\gamma} \mapsto x^{\gamma} u^{\langle w, \gamma\rangle} . \\
\Phi: f \mapsto \sum_{\gamma \in N \cap P} a_{\gamma} x^{\gamma} u^{\langle w, \gamma\rangle} .
\end{gathered}
$$

$\Phi f \in \mathbb{C}[N]$ might constrain some $a_{\gamma}$.
Akhtar, Coates, Galkin, Kasprzyk [arXiv:1212.1785]

## Motivation

## Definition (Kasprzyk, Tveiten)

Let $P$ be a canonical Fano polytope.
Consider

$$
f=\sum_{\gamma \in N \cap P} a_{\gamma} x^{\gamma}
$$

such that there exist $\Phi_{1}, \ldots, \Phi_{r}$ such that

$$
\Phi_{i} f \in \mathbb{C}[N] \quad \text { for } i=1, \ldots, r
$$

fixes uniquely all values of $a_{\gamma}, \quad \gamma \in \partial P \backslash \mathcal{V}(P)$,
then we call $f$ a (rigid) Maximally Mutable Laurent polynomial.

- if $\gamma \in \mathcal{V}(P)$, then $a_{\gamma}=1$.
- If $\gamma=(0, \ldots, 0)$, then $a_{\gamma}=0$.

Akhtar, Coates, Galkin, Kasprzyk [arXiv:1212.1785]

## Motivation

## Definition

Given a Laurent polynomial $f \in \mathbb{C}\left[x_{1}^{ \pm 1}, \ldots, x_{n}^{ \pm 1}\right]$, the period of $f$ is:
$\pi_{f}(t)=\frac{1}{(2 \pi \mathrm{i})^{n}} \int_{\left|x_{1}\right|=1, \cdots,\left|x_{n}\right|=1} \frac{1}{1-t f\left(x_{1}, \ldots, x_{n}\right)} \frac{\mathrm{d} x_{1}}{x_{1}} \cdots \frac{\mathrm{~d} x_{n}}{x_{n}}$,
Moreover, if $f$ is mirror to a Kano manifold,

$$
\pi_{f}(t)=\sum_{k \geq 0} c_{k} t^{k} \quad c_{k} \in \mathbb{Z}, \quad \text { for all } k
$$

The sequence $\left(c_{0}, c_{1}, \ldots\right)$ is called the period sequence of $f$.

$$
c_{k}=\operatorname{coeff}_{1} f^{k}
$$

[Akhtar, Coates, Galkin, Kasprzyk (arXiv:1212.1785)]

## Motivation

## example



$$
f=x+y+\frac{1}{x y}
$$

$$
\pi_{f}(t)=1+6 t^{3}+90 t^{6}+1680 t^{9}+34650 t^{12}+756756 t^{15}+\ldots
$$

## Motivation

example


$$
f=y+\frac{1}{x y}\left(x^{2}+a x+1\right)
$$

$f$ is Maximally Mutable iff $a=2$, and we have:

$$
\begin{gathered}
f=y+\frac{1}{x y}(x+1)^{2} \\
\pi_{f}(t)=1+4 t^{2}+36 t^{4}+400 t^{6}+4900 t^{8}+63504 t^{10}+\ldots
\end{gathered}
$$

## Motivation

If $\operatorname{dim} P>2, P$ may have more than one MM Laurent polynomials. example


$$
f=\frac{1}{z}+z\left(\frac{(x+1)^{2}}{x y}+\frac{x^{2}+2 x+1}{x}+y\right)
$$

$$
g=\frac{1}{z}+z\left(\frac{(x+1)^{2}}{x y}+\frac{x^{2}+3 x+1}{x}+y\right)
$$

(left)

$$
\pi_{f}(t)=1+4 t^{2}+60 t^{4}+1120 t^{6}+24220 t^{8}+\ldots
$$

(right)

$$
\pi_{g}(t)=1+6 t^{2}+90 t^{4}+1860 t^{6}+44730 t^{8}+\ldots
$$

## Computational aspects

## Use case

Given a class of polytopes $\mathcal{U}$, for all $P \in \mathcal{U}$ :

- find all Maximally-Mutable polynomials supported on $P$.
- Make a list of the period sequences.
- Compute other interesting quantities.
$\mathcal{U}$ can be very large

| 3-reflexives | 3-canonicals | 4-reflexives |
| :---: | :---: | :---: |
| 4319 | 674688 | 473800776 |

## Performance Requirements

- Ideally, as many cores as polytopes in $\mathcal{U}$.
- High single-core performance.
- High disk read/write rate.


## Computational aspects

IO infrastructure

- Text/binary file dump.
- As simple as possible.
- Naturally concurrent.
- Slow postprocessing.
- SQLite database.
- High quality, simple to use SQL database.
- Not designed for concurrency.
- Dedicated database server.
- High performance, high concurrency.
- Flexible data structures, SQL (Postgres, CockroachDB) or JSON (Mongo).
- Non negligible set up and maintenance cost.
- The HPC-database interface is of fundamental importance.

Our solution: in-house tooling for workers/database interaction [Coates, Kasprzyk (https://bitbucket.org/pcas)]

## Example: canonical Fano 3-topes

Canonical Fano polytopes in 3 dimensions were classified by Kasprzyk [arXiv:0806.2604].
Up to GL( $\mathbb{Z}, 3)$-equivalence, there are 674688 canonical Fano 3-topes.
Target For each canonical Fano 3-tope $P$ :

1. find all Maximally Mutable polynomials $f$ supported on $P$.
2. For every such $f$, compute the period sequence.
3. Can we analise and understand this data?

## Example: code metrics

The Maximally Mutables code is quite complex The most demanding computational tasks include:

- Minkowski decompositions.
- Gröbner basis computations.
- Graph search algorithm.
- Several performance tricks.

Performance analysis and optimisation

- Code profiling (development).
- Runtime metrics (deployment).


## Example: code metrics

Random sample of 10000 polytopes.


Running time by task


## Example: canonical Fano 3-topes

As of today, the search is $99.8 \%$ complete.
We found $\sim 170000$ Maximally Mutable Laurent polynomials.
These polynomials originate $\sim 8200$ period sequences.

- We can leverage the wealth of open source libraries for Data Science.
- Example: basic analysis of the period sequence data.
- Tools: pandas and scikit-learn libraries for python.


## Example: canonical Fano 3-topes

## Data Analysis

Difficulties in dealing with period sequences data:

- the coefficients of the period sequences grow quickly. $\hookrightarrow$ Work instead with $\log c_{k}$, then rescale so that $\log c_{k} \in[-1,1]$
- The data is high-dimensional (in this computation, each datapoint is in $\mathbb{Z}^{14}$ ( $[-1,1]^{14}$ after rescaling).
$\hookrightarrow$ Principal Component Analysis: project to linear subspaces that explain most of the variance in the data.
- Beware of low-dimensional projections of high-dimensional data.


## Example: canonical Fano 3-topes

Low-dimensional projection
Every point represents a period sequence.
Period sequences of smooth manifolds are in red.


## Example: canonical Fano 3-topes

Low-dimensional projection
Every point represents a period sequence.
Anticanonical degree.


## Example: canonical Fano 3-topes

Low-dimensional projection
Every point represents a period sequence.
Gorenstein index (log scale).


## Perspectives

Opportunities and challenges ahead

- Databases are evolving quickly.
$\hookrightarrow$ Automatic sharding (CockroachDB).
$\hookrightarrow$ Optimised performance for SSDs.
$\hookrightarrow$ Memory caches (e.g. Redis).
- Data analysis on $\sim 1 \mathrm{~Tb}$.

