# Conjugation in Groups of Lie type 

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## Groups of Lie type

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$G$ a reductive algebraic group over the field $K$.

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$G$ a reductive algebraic group over the field $K$.

- Fix a maximal torus $T$ of $G$,
- and let $\left(X, \Phi, Y, \Phi^{\star}\right)$ be its root datum,
- and W its Weyl group.


## Steinberg presentation

## Generators

- $x_{\alpha}(a)$, for $\alpha \in \Phi$ and $a \in K$,
- $\gamma \otimes t$, for $y \in Y$ and $t \in K^{*}$.


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- $n_{\alpha}=x_{\alpha}(\mathrm{I}) x_{-\alpha}(-\mathrm{I}) x_{\alpha}(\mathrm{I})$.


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Subgroups:

- T, generated by $\gamma \otimes t$,
- $N$, generated by $n_{\alpha}$,
- $U$, generated by $x_{\alpha}(a)$, with $\alpha \in \Phi^{+}$,
- $U^{-}$, generated by $x_{\alpha}(a)$, with $\alpha \in \Phi^{-}$,
- B=TU.


## Representations

Fix a rational representation

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\rho: G \rightarrow \mathrm{GL}(V)
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## Representations

Fix a rational representation

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which is a $G$-module by

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$$

Weights :

$$
V_{\mu}=\left\{v \in V \mid v(\gamma \otimes t)=t^{\langle\mu, \gamma\rangle} v \text { for all } \gamma \in Y, t \in K^{*}\right\}
$$

now $V=\bigoplus V_{\mu}$.

## Weight Base

Definitions:

- $n$ is the rank of the root system of $G$,
- $\delta: G \rightarrow G: x_{\alpha}(a) \mapsto x_{-\alpha}(-a)^{-1}$,
- $I \subseteq\{\mathrm{I}, \ldots, n\}: W_{I}=\left\langle n_{\alpha_{i}} \mid i \in I\right\rangle$,
- "Levi complement": $L_{I}=B W_{I} B \cap B^{\delta} W_{I}^{\delta} B^{\delta}$.


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## Calculating a weight base

(Cohen, Murray, Taylor, 2003)
In: Reductive algebraic group $G$ over $K$, Highest weight representation $\rho: G \rightarrow \mathrm{GL}(V)$,
Out: $\quad \lambda_{\mathrm{I}}, \ldots, \lambda_{k} \in V, J_{\mathrm{I}}, \ldots, J_{k} \subseteq\{\mathrm{I}, \ldots, n\}$, such that $G=L_{J_{\mathrm{I}}} \geq L_{J_{2}} \geq \ldots \geq L_{J_{k+1}}=T$, and $\lambda_{i}$ is a highest weight for $L_{J_{i}}$ acting on $V$.

## Original

## Row and column reduction

(Cohen, Murray, Taylor, 2003)
In: Reductive algebraic group $G$ over $K$, Representation $\rho: G \rightarrow \mathrm{GL}(V)$, $A \in \rho(G)$.
Out: $\quad g \in G$ such that $A=\rho(g)$.
Algorithm: A generalization of row and column reduction.

## For vectors (I)

## Row reduction for vectors (I)

In: Reductive algebraic group $G$ over $K$, Representation $\rho: G \rightarrow \mathrm{GL}(V)$, $\left\{z_{\mathrm{I}}, \ldots, z_{q}\right\} \subseteq V$,
Out: $\quad g \in G$ such that, for all $i \in\{\mathrm{I}, \ldots, q\}$ :

$$
z_{i} \mathrm{~g}^{-\mathrm{I}} \in \bigoplus_{\mu \succeq \lambda_{i}} V_{\mu}
$$

## Into Borel

## Conjugation into the Borel subgroup

In: Reductive algebraic group $G$ over $K$, $c$ a semisimple element of $G$,
Out: $\quad d \in B, a \in G$ such that $d=c^{a}$

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In: Reductive algebraic group $G$ over $K$, $c$ a semisimple element of $G$,
Out: $d \in B, a \in G$ such that $d=c^{a}$
Sketch of the algorithm:

1. Compute centraliser $\mathfrak{c}$ of $c$ in the Lie algebra $\mathfrak{g}$,
2. Compute splitting Cartan subalgebra $\mathfrak{h}$ contained in $\mathfrak{c}$,
3. Write down basis of $\mathfrak{g}$ with respect to $\mathfrak{h}$,
4. Row reduce some entries of this basis, obtaining $g \in G$,
5. $\mathrm{g}^{-\mathrm{I}}$ does the job.

## Into Borel: Example

We consider $A_{3}$ over $\mathbb{F}_{59}$.

$$
c=\left(\begin{array}{cccc}
9 & \mathrm{I} & 53 & 8 \\
42 & 0 & 8 & 44 \\
48 & 0 & 7 & 23 \\
\mathrm{I} 2 & \mathrm{I} & 23 & 8
\end{array}\right)
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- Centraliser is 5 dimensional over $\mathbb{F}_{59}$,
- Cartan Subalgebra is 3 dimensional over $\mathbb{F}_{59}^{4}$,
- We row reduce:

$$
\left(\begin{array}{ccccccccccccccc}
4 & 25 & 32 & 18 & 23 & 5 & 35 & 39 & 9 & 54 & 52 & 44 & 20 & 29 & 52 \\
55 & 34 & 0 & 4 \mathrm{I} & 0 & 2 & 49 & 17 & 24 & 30 & 9 & 0 & 57 & 47 & 42
\end{array}\right)
$$

## Into Borel: Example

We consider $A_{3}$ over $\mathbb{F}_{59}$.

- And find

$$
\mathrm{g}^{-\mathrm{I}}=\left(\begin{array}{cccc}
0 & 0 & 0 & 58 \\
0 & 0 & \mathrm{I} & 2 \mathrm{I} \\
0 & 58 & 50 & 25 \\
\mathrm{I} & 5 \mathrm{I} & 30 & 53
\end{array}\right)
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\end{array}\right)
$$

- Now

$$
c^{\prime}=g c g^{-\mathrm{I}}=\left(\begin{array}{cccc}
25 & 0 & 0 & 0 \\
7 & 28 & 0 & 0 \\
24 & 44 & 50 & 0 \\
8 & \mathrm{I} & 0 & 50
\end{array}\right)
$$

## Towards the standard torus

$$
c^{\prime}=\left(\begin{array}{cccc}
25 & 0 & 0 & 0 \\
7 & 28 & 0 & 0 \\
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- Recall $\delta: G \rightarrow G: x_{\alpha}(a) \mapsto x_{-\alpha}(-a)^{-1}$,


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- Recall $\delta: G \rightarrow G: x_{\alpha}(a) \mapsto x_{-\alpha}(-a)^{-1}$,
- Now

$$
c^{\prime \delta}=\left(\begin{array}{cccc}
26 & 36 & 3 \mathrm{I} & 53 \\
0 & \mathrm{I} 9 & \mathrm{I} 2 & 48 \\
0 & 0 & \mathrm{I} 3 & 0 \\
0 & 0 & 0 & 13
\end{array}\right)
$$

- So we would like to do the same algorithm again...


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$$

- So we would like to do the same algorithm again...
- but that will not work.


## For vectors (II)

## Row reduction for vectors (II)

In: Reductive algebraic group $G$ over $K$, Representation $\rho: G \rightarrow \mathrm{GL}(V)$, $\left\{z_{\mathrm{I}}, \ldots, z_{q}\right\} \subseteq V$,
Out: $\quad g \in G$ such that, for all $i \in\{\mathrm{I}, \ldots, q\}$ :

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Out: $\quad g \in G$ such that, for all $i \in\{\mathrm{I}, \ldots, q\}$ :

$$
\begin{gathered}
z_{i} \mathrm{~g}^{-\mathrm{I}} \in \bigoplus_{\mu \succeq \lambda_{i}} V_{\mu} \\
\text { and, for } i \in\{\mathrm{I}, \ldots, k\}, j \in\{\mathrm{I}, \ldots, i-\mathrm{I}\}, \\
z_{i} \cdot v_{i} \neq \mathrm{o} \text { and } z_{i} \cdot v_{j}=\mathrm{o} \Rightarrow \mathrm{~g} \in B^{\delta} .
\end{gathered}
$$

## Into standard torus

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In: Reductive algebraic group $G$ over $K$, $c$ a semisimple element of $G$,
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In: Reductive algebraic group $G$ over $K$, $c$ a semisimple element of $G$,
Out: $\quad h \in T, a \in G$ such that $h=c^{a}$
Sketch of the algorithm:

1. Compute $c^{\prime} \in B$,
2. Conjugate $c^{\prime \delta}$ into $B$ as before, but use new row reduction,
3. Compose results.

## Into standard torus: Example

- Now

$$
c^{\prime \delta}=\left(\begin{array}{cccc}
26 & 36 & 3 \mathrm{I} & 53 \\
0 & \text { I9 } & \text { I2 } & 48 \\
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c^{\prime \delta}=\left(\begin{array}{cccc}
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\end{array}\right)
$$

- We row reduce:

$$
\left(\begin{array}{ccccccccccccccc}
\mathrm{I} & \mathrm{I} 4 & \mathrm{O} & \mathrm{O} & 0 & 45 & 34 & \mathrm{I} 9 & \mathrm{I} 3 & 25 & 0 & 0 & 0 & 54 & \mathrm{I} 9 \\
9 & 8 & 42 & 0 & 57 & 38 & \mathrm{I} & 58 & 58 & 0 & 0 & \mathrm{I} 5 & 0 & 22 & 0
\end{array}\right)
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- Now

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9 & 8 & 42 & 0 & 57 & 38 & \mathrm{I} & 58 & 58 & 0 & 0 & \mathrm{I} 5 & 0 & 22 & 0
\end{array}\right)
$$

- And obtain $g^{\prime}$, such that

$$
g^{\prime} c^{1 \delta} g^{\prime-1}=\left(\begin{array}{cccc}
33 & \circ & 0 & 0 \\
0 & 40 & 0 & 0 \\
0 & \circ & 46 & 0 \\
0 & \circ & 0 & 46
\end{array}\right)
$$

## Weighted Dynkin Diagrams

## "Definition"

In: Lie algebra $\mathfrak{g}$, with simple root system $\Delta$, $e \in \mathfrak{g}$ nonzero nilpotent,
Out: $\quad p_{\alpha} \in\{0, \mathbf{I}, 2\}, \alpha \in \Delta$.

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## Why weighted Dynkin diagrams are useful

Theorem: $e$ and $e^{\prime}$ have the same weighted Dynkin diagram if and only if they are in the same $G$-orbit.

## Computing Weighted Dynkin Diagrams

Let $e \in \mathfrak{g}$, and $\langle e, h, f\rangle$ the $\mathfrak{s l}_{2}$ subalgebra containing $e$.

## Computing

Computing the weighted Dynkin diagram for $e$ is easy if $h$ in the standard torus t .

## TU/e

## Computing Weighted Dynkin Diagrams

## Example in MAGMA (I)

> G := GroupOfLieType("D4", Rationals());
> rho,L : = AdjointRepresentation (G);
> rhoL := AdjointRepresentation(L);
$>$
$>$ e1 := L.1;
> lbl1 := WeightedDynkinDiagramLabels(e1);
> lbl1;
$[0,1,0,0]$

## Computing Weighted Dynkin Diagrams

## Example in MAGMA (II)

> lbl1;
[ 0, 1, 0, 0 ]
$>\mathrm{c}:=\mathrm{elt}\langle\mathrm{G}| \mathrm{[ }$ ( <1,3>, <2,3>, 2, 3 *] >; c;
x2(3) x5(9) x1(3) n2 n3
> C := ChangeRing(rho(c), BaseRing(L));
$>$
$>$ e2 := L! (e1*C);
> e2;
(-3 00010 ....... 0000$)$
$>$
> lbl2 := WeightedDynkinDiagramLabels(e2);
> lbl2;
$[0,1,0,0$ ]

## Conclusion

- New variants of row reduction,
- Application: Conjugation,
- Application: Weighted Dynkin diagrams.


## Future Research

- Small characteristic cases,
- Fix bugs,
- Find conjugators for subalgebras in the same orbit,


## Questions?

